Exercises (Context-Free Languages)

1 Context-Free Grammars

Exercise: Give context-free grammars that generate the following languages.

1. $L := \{a^k b^l c^{k+l} \mid k, l \in \mathbb{N}\}$ 2. $L := \{a^k b^k c^l d^l \mid k, l \in \mathbb{N}\}$ 3. $L := \{a^k b^l \mid k \ge 1, l > k\}$ 4. $L := \{w \in \{a, b\}^* \mid |w| \text{ odd}, a \text{ in middle}\} (= \{uav \in \{a, b\}^* \mid |u| = |v|\})$

Solution:

- 1. $S \to aSc \mid B$ $B \to bBc \mid \varepsilon$
- 2. $S \rightarrow AC$ $A \rightarrow aAb \mid \varepsilon$ $C \rightarrow cCd \mid \varepsilon$
- 3. $S \rightarrow AB$ $A \rightarrow aAb \mid ab$ $B \rightarrow bB \mid b$
- 4. $S \rightarrow aSa \mid aSb \mid bSa \mid bSb \mid a$ or shorter: $S \rightarrow ASA \mid a$ $A \rightarrow a \mid b$

Exercise: Show that context-free languages are closed under the reversal operation. **Solution:** Idea: reverse all right-hand sides of productions

2 Context-Free and Regular Languages

Exercise: Show that every regular expression can directly be translated into an equivalent context-free grammar.

Solution: Given regular expression α over Σ , we inductively construct CFG $G_{\alpha} = \langle N, \Sigma, P, S \rangle$ as follows:

- $\alpha = \emptyset$: $P = \emptyset$
- $\alpha = \varepsilon$: $P = \{S \to \varepsilon\}$
- $\alpha = a$: $P = \{S \rightarrow a\}$
- $\alpha = \alpha_1 + \alpha_2$: $P = \{S \to S_1 \mid S_2\} \cup P_1 \cup P_2 \text{ (where } G_{\alpha_i} = \langle N_i, \Sigma, P_i, S_i \rangle \text{ for } i = 1, 2)$
- $\alpha = \alpha_1 \cdot \alpha_2$: $P = \{S \to S_1 S_2\} \cup P_1 \cup P_2 \text{ (where } G_{\alpha_i} \text{ as above)}$
- $\alpha = \alpha_1^*$: $P = \{S \to S_1 S \mid \varepsilon\} \cup P_1 \text{ (where } G_{\alpha_1} \text{ as above)}$