

## Verifying Al-controlled Hybrid Systems Master's Thesis Presentation

Gabriela Jiang

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- Hybrid systems: systems combining continuous and discrete dynamics
- If too complex, may be regulated by neural network controllers
- Often in engineering and control systems

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#### Examples

Automotive systems, robotics, power systems

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- Often in engineering and control systems

#### Examples

Automotive systems, robotics, power systems

#### Importance

- Ensure safety and reliability in critical systems
- Predict and prevent undesirable behaviors

## Challenges

- Complexity of modeling and analyzing hybrid systems
- Black box behavior of neural networks
- Dealing with over-approximation

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### Current methods

- Model checking
- Theorem proving
- Simulation-based approaches

## Challenges

- Complexity of modeling and analyzing hybrid systems
- Black box behavior of neural networks
- Dealing with over-approximation
- Current methods
  - Model checking
  - Theorem proving
  - Simulation-based approaches
- Limitations of current methods
  - Limited scalability
  - Assumptions about system behavior
  - Lack of guarantees for certain properties

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Autonomous system

$$\dot{x}(t) = Ax(t)$$

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Non-autonomous system

$$\dot{x}(t) = Ax(t) + \frac{Bu(t)}{Bu(t)}$$

where x are the variables, A is the flow matrix at time t, u the external inputs.

 $(X, U, F, \mathcal{N}, \delta, \mathcal{X}_0)$ 

 $(\mathbf{X}, U, F, \mathcal{N}, \delta, \mathcal{X}_0)$ 

$$\begin{array}{c|c}
 & \text{NN Controller} \\
 & u(t_k) = \mathcal{N}(y(t_k)) \\
\hline & y(t_k) = Cx(t_k) \\
\hline & \dot{x}(t) = Ax(t) + Bu(t_k) \\
\hline & \dot{x}(t)$$

 $X = \{x_1, \ldots, x_n\}$  is the set of finitely many state variables

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\end{array}$$

 $U = \{u_1, \ldots, u_m\}$  the set of finitely many control variables

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$$\begin{array}{c|c}
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### $F: \mathbb{R}_{\geq 0} \to \mathbb{R}^n$ the ODE defining the continuous dynamics of the plant

 $(X, U, F, \mathcal{N}, \delta, \mathcal{X}_0)$ 

$$(1)$$
NN Controller
$$u(t_k) = \mathcal{N}(y(t_k))$$

$$(1)$$

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$$(1)$$

$$V(t_k) = Cx(t_k)$$

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### $\mathcal{N}: \mathbb{R}^{n_{x}} \to \mathbb{R}^{n_{y}}$ denotes the input-output mapping by the NN controller

 $(X, U, F, \mathcal{N}, \frac{\delta}{\delta}, \mathcal{X}_0)$ 

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### $\delta \in \mathbb{R}_{>0}$ denotes the control step size

 $(X, U, F, \mathcal{N}, \delta, \mathcal{X}_0)$ 

### $\mathcal{X}_0 \subset \mathbb{R}^n$ is the set of the initial values of the state variables

### Modeled as

Neural network

Simplified hybrid automaton

#### Reachability analysis

Exact star-based NN reachability

Flowpipe construction



















$$\mathcal{R}(I) = \{ \sigma \in \Sigma \mid \exists \sigma' \in I.\sigma' \to^* \sigma \}$$

not computable [Henzinger et al. 1998]

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Hence over-approximate approaches where

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Such as flowpipe construction

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Hence over-approximate approaches where

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- Such as flowpipe construction
- In the following,  $T = K\delta$

# **Flowpipe Construction**





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$$X_0 \cup e^{\delta A} X_0$$

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# **Flowpipe Construction**



$$conv(X_0 \cup e^{\delta A}X_0)$$

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# Flowpipe Construction



$$\Omega_0 = \mathit{conv}(X_0 \cup e^{\delta A}X_0) \oplus \mathcal{B}_lpha$$

# Flowpipe Construction



$$egin{aligned} \Omega_0 &= \mathit{conv}(X_0 \cup e^{\delta A}X_0) \oplus \mathcal{B}_lpha \ \Omega_{i+1} &= e^{\delta A}\Omega_i \end{aligned}$$

## State Set Representations







(c) V-polytope

## Reachability Analysis of Neural Networks

### Definition (Reachable set of an FNN)

Given

- Input set  $I = \{x \in \mathbb{R}^d \mid Ax \leq b\}$  with  $A \in \mathbb{R}^{n \times d}, b \in \mathbb{R}^n$
- *k*-layer FNN  $\mathcal{N} = \{L_1, \cdots, L_k\}$

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• Reachable set  $\mathcal{N}(I) = \mathcal{R}_{L_k}$  is defined by

$$\begin{aligned} \mathcal{R}_{L_0} &:= I \\ \mathcal{R}_{L_i} &:= \left\{ y_i \mid y_i = \phi_i \left( \mathcal{W}_i y_{i-1} + b_i \right), y_{i-1} \in \mathcal{R}_{L_{i-1}} \right\}, \quad i = 1, \dots, k \end{aligned}$$

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### Approach [Tran 2020]

Exact analysis for FNNs with piece-wise linear activations.

A generalized star set is a tuple  $\Theta = \langle c, V, P \rangle$  where

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Efficient w.r.t. computing affine mappings and halfspace intersections.

## Example Starset





 $\Theta + t$ 

Translate by

$$t = \left[ egin{array}{c} -1 \\ -1 \end{array} 
ight]$$



Rotate by

$$M = \begin{bmatrix} \cos 45^{\circ} & -\sin 45^{\circ} \\ \sin 45^{\circ} & \cos 45^{\circ} \end{bmatrix}$$

# Example Starset



$$c' = \begin{bmatrix} -1\\ -1 \end{bmatrix},$$

$$V' = \left\{ \begin{bmatrix} \cos 45^{\circ}\\ \sin 45^{\circ} \end{bmatrix}, \begin{bmatrix} -\sin 45^{\circ}\\ \cos 45^{\circ} \end{bmatrix} \right\},$$

$$P = \begin{bmatrix} -1 & 0\\ 1 & 0\\ 0 & -1\\ 0 & 1 \end{bmatrix} \begin{bmatrix} x\\ y \end{bmatrix} \le \begin{bmatrix} -4\\ 4\\ -3\\ 3 \end{bmatrix}$$

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Develop techniques for verifying the safety of NNCS

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- Focus on non-autonomous hybrid systems with discrete interactions

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- Focus on non-autonomous hybrid systems with discrete interactions
- Keep over-approximation errors small

### 1 Continuous

 $\blacksquare$  Independent of the time step  $\delta$ 

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- $\blacksquare$  Independent of the time step  $\delta$
- 2 Discrete
  - Irregular: time steps  $\delta_1, \delta_2, \ldots$
  - $\blacksquare$  Regular: time step  $\delta$



(a)  $\delta = 0.01$ 







Trade-off between safety and computational effort



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- Devise a suitable controller to ensure the system's safety
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  - Train controller to learn behavior of hybrid automata
  - Reduce hybrid automaton to a simplified automaton

## Simplified Hybrid Automaton

$$\mathcal{H}' = (Loc', Var', Lab', Act', Init')$$

$$k \in [19.5, 20.5], f_{2,x} \rightarrow \begin{pmatrix} \ell \\ f_{1,x} : \dot{x} = -Kx \\ f_{2,x} : \dot{x} = -K(x-\rho) \end{pmatrix}$$

## Simplified Hybrid Automaton

$$\mathcal{H}' = (Loc', Var', Lab', Act', Init')$$

$$\ell$$

$$f_{1,x} : \dot{x} = -Kx$$

$$f_{2,x} : \dot{x} = -K(x - \rho)$$
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- 2 Discrete value from a finite and countable set
  - Multi-class classifier
  - Variable in the hybrid system's flow function

$$\dot{x}(t) = Ax(t) + Bu(t)$$

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Flows of the thermostat:

$$\begin{aligned} \dot{x} &= -Kx \qquad (\ell_{\text{off}}) \\ \dot{x} &= K(\rho - x) = -Kx + K\rho \qquad (\ell_{\text{on}}) \end{aligned}$$

where  $K\rho$  can be seen as the external input Bu(t) from the controller

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where  $K\rho$  can be seen as the external input Bu(t) from the controller • Hence, let the neural network predict  $K\rho$ 

• Or more generally,  $h_i$ 

Use step function

$$f_{\text{step}}(x) = \sum_{i=1}^{n} \alpha_i \chi_{A_i}(x)$$

where  $\alpha_i \in \mathbb{R}$ , n > 0, and

$$\chi_{A_i}(x) = egin{cases} 1, \ ext{if } x \in A_i, \ 0, \ ext{otherwise} \end{cases}$$

for each interval  $A_i$ ,  $i = 1, \ldots, n$ .

Given: ordered target outputs h<sub>i</sub> ∈ ℝ with i = 1,..., n and h<sub>i-1</sub> < h<sub>i</sub>
Apply step function with intervals

$$A_{i} = \begin{cases} (-\infty, \frac{h_{i}+h_{i+1}}{2}), & \text{if } i = 1, \\ [\frac{h_{i}+h_{i-1}}{2}, \infty), & \text{if } i = n, \\ [\frac{h_{i-1}+h_{i}}{2}, \frac{h_{i}+h_{i+1}}{2}), & \text{otherwise} \end{cases}$$

and  $\alpha_i = h_i$ .

#### Output Classification: Visualized



#### Neural Network Architecture





NN reachability analysis

where  $x_j \subset \mathbb{R}$ with  $\Theta_{out}$  covering trivial sets  $\{h_i\}$ 

Example (Thermostat)

At time point  $t_k$ , let

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Example (Thermostat)

At time point  $t_k$ , let

- temperature  $x \in [x_{k_l}, x_{k_u}]$ ,
- mode  $m_k \in \{0, 1\}$  (off / on).

Define  $\Theta_{I_k} = \langle c, V, P \rangle$  with  $P(\alpha) \triangleq C\alpha \leq s$  s.t.

$$c = \begin{bmatrix} 0 \\ 0 \end{bmatrix}, \quad V = \{e_1, e_2\}, \quad C = \begin{bmatrix} -1 & 0 \\ 1 & 0 \\ 0 & -1 \\ 0 & 1 \end{bmatrix}, \quad s = \begin{bmatrix} -x_{k_l} \\ x_{k_u} \\ -m_k \\ m_k \end{bmatrix}$$

#### Definition (Initial set of NNCS analysis)

- Let  $(X, U, F, \mathcal{N}, \delta, \mathcal{X}_0)$  be an NNCS and  $T = K\delta$  the time horizon
- Initial set  $\mathcal{I}_k$  at time point  $t_k = k\delta$  is defined by

$$egin{aligned} \mathcal{I}_0 &:= \mathcal{X}_0 \ \mathcal{I}_k &:= \mathsf{e}^{\delta A_{k-1}} \mathcal{I}_{k-1}, \quad k = 1, \dots, K \end{aligned}$$

where  $A_{k-1}$  is the flow matrix over the time interval  $[t_{k-1}, t_k]$ .

# Handling Potential Branching



 $\mathcal{S}_{i-1} = \{ (\mathcal{X}_i, F, \mathcal{I}) \mid \mathcal{X}_i \subseteq V', F \in Act'(\ell), \mathcal{I} \in \textit{Init'} \}$ 

# Handling Potential Branching



 $\mathcal{S}_{i-1} = \{(\mathcal{X}_i, f_{1,x}, \mathcal{I}_{i-1}), (\mathcal{X}'_i, f_{2,x}, \mathcal{I}_{i-1})\}$ 

Algorithm NNCS analysis (discrete).

**Input:** NNCS  $(X, U, F, \mathcal{N}, \delta, \mathcal{X}_0)$ , HA  $\mathcal{H}'$ , initial set  $\mathcal{I}_0, T = K\delta$ **Output:** Reachable states  $\mathcal{R}$  over [0, T]1: procedure NNCSREACH(I\_0, K)  $\mathcal{R} \leftarrow \text{CONSTRUCTFLOWPIPE}(\mathcal{I}_0, \delta);$ 2.  $\mathcal{X}_i \leftarrow \text{GETRELEVANTSETS}(\mathcal{R}, \delta);$ 3. 4:  $S_{i-1} \leftarrow \{(\mathcal{X}_i, F_0, \mathcal{I}_0)\}$ for i = 1 to K do 5:  $S_i \leftarrow \emptyset$ : 6: for  $(\mathcal{X}_i, F_i, \mathcal{I}_i) \in \mathcal{S}_{i-1}$  do 7:  $(\mathcal{X}', \mathcal{F}', \mathcal{I}') \leftarrow \text{COMPUTESETFLOWMAPS}(\mathcal{X}_i, \mathcal{F}_i, \mathcal{I}_i);$ 8:  $S_i \leftarrow S_i \cup \{(\mathcal{X}', F', \mathcal{I}')\}$ : Q٠ 10: end for for j = 1 to  $|S_j|$  do 11:  $\mathcal{R}_i \leftarrow \text{CONSTRUCTFLOWPIPE}(\mathcal{I}_i, \delta);$ 12:  $\mathcal{R} \leftarrow \mathcal{R} \cup \mathcal{R}_i$ 13: end for 14.  $S_{i-1} \leftarrow S_i$ ; 15 16.  $\mathcal{X}_i \leftarrow \text{GETRELEVANTSETS}(\mathcal{R}, (i+1)\delta);$ 17: end for return  $\mathcal{R}$ 18:

19: end procedure

 $t_0 = 0$ 



$$\mathcal{I}_0 = \left\{ (f_{2,x}, \nu) \in \Sigma' \mid \nu(x) \in [19.5, 20.5] \right\}$$

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**Input:** NNCS  $(X, U, F, \mathcal{N}, \delta, \mathcal{X}_0)$ , HA  $\mathcal{H}'$ , initial set  $\mathcal{I}_0, T = K\delta$ **Output:** Reachable states  $\mathcal{R}$  over [0, T]1: procedure NNCSREACH( $\mathcal{I}_0, K$ )  $\mathcal{R} \leftarrow \text{CONSTRUCTFLOWPIPE}(\mathcal{I}_0, \delta);$ 2:  $\mathcal{X}_i \leftarrow \text{GETRELEVANTSETS}(\mathcal{R}, \delta);$ 3:  $S_{i-1} \leftarrow \{(\mathcal{X}_i, F_0, \mathcal{I}_0)\}$ 4: 5. for i = 1 to K do  $S_i \leftarrow \emptyset$ : 6. for  $(\mathcal{X}_i, F_i, \mathcal{I}_i) \in \mathcal{S}_{i-1}$  do 7:  $(\mathcal{X}', F', \mathcal{I}') \leftarrow \text{COMPUTESETFLOWMAPS}(\mathcal{X}_i, F_i, \mathcal{I}_i);$ 8: 9:  $S_i \leftarrow S_i \cup \{(\mathcal{X}', F', \mathcal{I}')\};$ 10: end for 11. for i = 1 to  $|S_i|$  do  $\mathcal{R}_i \leftarrow \text{CONSTRUCTFLOWPIPE}(\mathcal{I}_i, \delta);$ 12.  $\mathcal{R} \leftarrow \mathcal{R} \cup \mathcal{R}_i;$ 13. end for 14:  $S_{i-1} \leftarrow S_i$ ; 15:  $\mathcal{X}_i \leftarrow \text{GETRELEVANTSETS}(\mathcal{R}, (i+1)\delta);$ 16: 17: end for return  $\mathcal{R}$ 18. 19: end procedure

 $t_0 = 0$ 



Flowpipe construction for  $\delta$  time

 $\mathcal{R} = \Omega_{[0,\delta]}(\mathcal{X}_0, \mathcal{U})$ 

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 $t_0 = 0$ 



$$\mathcal{X}_i = \mathcal{R} \cap \{ \nu \mid \nu(t) = \delta \}$$
  
=  $\{ \nu \mid \nu(t) \in \delta, \nu(x) \in [20, 21] \}$ 

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 $t_1 = \delta$ 



$$\begin{aligned} \mathcal{X}_i &= \mathcal{R} \cap \{ \nu \mid \nu(t) = \delta \} \\ &= \{ \nu \mid \nu(t) \in \delta, \nu(x) \in [20, 21] \} \\ \mathcal{S}_{i-1} &= \{ (\mathcal{X}_i, f_{2,x}, \mathcal{I}_0) \} \end{aligned}$$

Algorithm NNCS analysis (discrete).

**Input:** NNCS  $(X, U, F, \mathcal{N}, \delta, \mathcal{X}_0)$ , HA  $\mathcal{H}'$ , initial set  $\mathcal{I}_0$ ,  $T = K\delta$ **Output:** Reachable states  $\mathcal{R}$  over [0, T]1: procedure NNCSREACH(I\_0, K)  $\mathcal{R} \leftarrow \text{CONSTRUCTFLOWPIPE}(\mathcal{I}_0, \delta);$ 2.  $\mathcal{X}_i \leftarrow \text{GETRELEVANTSETS}(\mathcal{R}, \delta);$ 3. 4:  $S_{i-1} \leftarrow \{(\mathcal{X}_i, F_0, \mathcal{I}_0)\}$ for i = 1 to K do 5:  $S_i \leftarrow \emptyset$ : 6: for  $(\mathcal{X}_i, F_i, \mathcal{I}_i) \in \mathcal{S}_{i-1}$  do 7:  $(\mathcal{X}', \mathcal{F}', \mathcal{I}') \leftarrow \text{COMPUTESETFLOWMAPS}(\mathcal{X}_i, \mathcal{F}_i, \mathcal{I}_i);$ 8:  $S_i \leftarrow S_i \cup \{(\mathcal{X}', F', \mathcal{I}')\}$ : Q٠ 10: end for for j = 1 to  $|S_j|$  do 11:  $\mathcal{R}_i \leftarrow \text{CONSTRUCTFLOWPIPE}(\mathcal{I}_i, \delta);$ 12:  $\mathcal{R} \leftarrow \mathcal{R} \cup \mathcal{R}_i$ 13: end for 14.  $S_{i-1} \leftarrow S_i$ ; 15 16.  $\mathcal{X}_i \leftarrow \text{GETRELEVANTSETS}(\mathcal{R}, (i+1)\delta);$ 17: end for return  $\mathcal{R}$ 18: 19: end procedure

 $t_1 = \delta$ 



$$\mathcal{X}_i = \mathcal{R} \cap \{ \nu \mid \nu(t) = \delta \}$$
  
= { $\nu \mid \nu(t) \in \delta, \nu(x) \in [20, 21]$ }  
 $S_{i-1} = \{ (\mathcal{X}_i, f_{2,x}, \mathcal{I}_0) \}$ 

Algorithm NNCS analysis (discrete).

**Input:** NNCS  $(X, U, F, \mathcal{N}, \delta, \mathcal{X}_0)$ , HA  $\mathcal{H}'$ , initial set  $\mathcal{I}_0, T = K\delta$ **Output:** Reachable states  $\mathcal{R}$  over [0, T]1: procedure NNCSREACH( $\mathcal{I}_0, K$ )  $\mathcal{R} \leftarrow \text{CONSTRUCTFLOWPIPE}(\mathcal{I}_0, \delta)$ : 2.  $\mathcal{X}_i \leftarrow \text{GETRELEVANTSETS}(\mathcal{R}, \delta);$ 3. 4:  $S_{i-1} \leftarrow \{(\mathcal{X}_i, F_0, \mathcal{I}_0)\}$ for i = 1 to K do 5:  $S_i \leftarrow \emptyset$ : 6: for  $(\mathcal{X}_i, F_i, \mathcal{I}_i) \in \mathcal{S}_{i-1}$  do 7:  $(\mathcal{X}', F', \mathcal{I}') \leftarrow \text{COMPUTESETFLOWMAPS}(\mathcal{X}_i, F_i, \mathcal{I}_i);$ 8.  $S_i \leftarrow S_i \cup \{(\mathcal{X}', F', \mathcal{I}')\};$ Q٠ 10: end for for i = 1 to  $|S_i|$  do 11:  $\mathcal{R}_i \leftarrow \text{CONSTRUCTFLOWPIPE}(\mathcal{I}_i, \delta);$ 12:  $\mathcal{R} \leftarrow \mathcal{R} \cup \mathcal{R}_i$ 13: end for 14.  $S_{i-1} \leftarrow S_i$ ; 15 16.  $\mathcal{X}_i \leftarrow \text{GETRELEVANTSETS}(\mathcal{R}, (i+1)\delta);$ end for 17: return  $\mathcal{R}$ 18: 19: end procedure

 $t_1 = \delta$ 



 $\mathcal{N}(\Theta(\mathcal{X}_{j})) = \Theta_{y} \text{ where } \Theta_{y}. \text{range} = \{K\rho\}$   $\implies F' = f_{2,x}$   $\implies \mathcal{I}' = \text{AFFINETRANSFORMATION}(\mathcal{I}_{0}, f_{2,x}, \delta)$   $\implies \mathcal{S}_{i} = \{(\mathcal{X}_{i}, f_{2,x}, \mathcal{I}')\}$ 

Algorithm NNCS analysis (discrete).

Input: NNCS  $(X, U, F, \mathcal{N}, \delta, \mathcal{X}_0)$ , HA  $\mathcal{H}'$ , initial set  $\mathcal{I}_0, T = K\delta$ **Output:** Reachable states  $\mathcal{R}$  over [0, T]1: procedure NNCSREACH(I\_0, K)  $\mathcal{R} \leftarrow \text{CONSTRUCTFLOWPIPE}(\mathcal{I}_0, \delta)$ : 2.  $\mathcal{X}_i \leftarrow \text{GETRELEVANTSETS}(\mathcal{R}, \delta);$ 3.  $S_{i-1} \leftarrow \{(\mathcal{X}_i, F_0, \mathcal{I}_0)\}$ 4: 5. for i = 1 to K do 6:  $S_i \leftarrow \emptyset_i$ for  $(\mathcal{X}_i, \mathcal{F}_i, \mathcal{I}_i) \in \mathcal{S}_{i-1}$  do 7:  $(\mathcal{X}', \mathcal{F}', \mathcal{I}') \leftarrow \text{COMPUTESETFLOWMAPS}(\mathcal{X}_i, \mathcal{F}_i, \mathcal{I}_i);$ 8:  $S_i \leftarrow S_i \cup \{(\mathcal{X}', F', \mathcal{I}')\}$ : Q٠ end for 10. for i = 1 to  $|S_i|$  do 11-12:  $\mathcal{R}_i \leftarrow \text{CONSTRUCTFLOWPIPE}(\mathcal{I}_i, \delta);$  $\mathcal{R} \leftarrow \mathcal{R} \cup \mathcal{R}_i$ 13: end for 14:  $S_{i-1} \leftarrow S_i$ : 15:  $\mathcal{X}_i \leftarrow \text{GETRELEVANTSETS}(\mathcal{R}, (i+1)\delta);$ 16 17. end for return  $\mathcal{R}$ 18: 19: end procedure

 $t_1 = \delta$ 



$$\mathcal{R}_j = \Omega_{[\delta, 2\delta]}(\mathcal{X}_j, \mathcal{U})$$
  
 $\mathcal{R} = \Omega_{[0, 2\delta]}(\mathcal{X}_0, \mathcal{U})$ 

Algorithm NNCS analysis (discrete).

**Input:** NNCS  $(X, U, F, \mathcal{N}, \delta, \mathcal{X}_0)$ , HA  $\mathcal{H}'$ , initial set  $\mathcal{I}_0$ ,  $T = K\delta$ **Output:** Reachable states  $\mathcal{R}$  over [0, T]1: procedure NNCSREACH(I\_0, K)  $\mathcal{R} \leftarrow \text{CONSTRUCTFLOWPIPE}(\mathcal{I}_0, \delta);$ 2.  $\mathcal{X}_i \leftarrow \text{GETRELEVANTSETS}(\mathcal{R}, \delta);$ 3. 4:  $S_{i-1} \leftarrow \{(\mathcal{X}_i, F_0, \mathcal{I}_0)\}$ for i = 1 to K do 5:  $S_i \leftarrow \emptyset$ : 6: for  $(\mathcal{X}_i, F_i, \mathcal{I}_i) \in \mathcal{S}_{i-1}$  do 7:  $(\mathcal{X}', \mathcal{F}', \mathcal{I}') \leftarrow \text{COMPUTESETFLOWMAPS}(\mathcal{X}_i, \mathcal{F}_i, \mathcal{I}_i);$ 8:  $S_i \leftarrow S_i \cup \{(\mathcal{X}', F', \mathcal{I}')\}$ : Q٠ 10: end for for j = 1 to  $|S_j|$  do 11:  $\mathcal{R}_i \leftarrow \text{CONSTRUCTFLOWPIPE}(\mathcal{I}_i, \delta);$ 12:  $\mathcal{R} \leftarrow \mathcal{R} \cup \mathcal{R}_i$ 13: end for 14.  $S_{i-1} \leftarrow S_i$ : 15 16  $\mathcal{X}_i \leftarrow \text{GETRELEVANTSETS}(\mathcal{R}, (i+1)\delta);$ 17: end for return  $\mathcal{R}$ 18: 19: end procedure

 $t_1 = \delta$ 



$$\begin{aligned} \mathcal{X}_i &= \mathcal{R} \cap \{ \nu \mid \nu(t) = 2\delta \} \\ &= \{ \nu \mid \nu(t) \in \delta, \nu(x) \in [20.5, 21.5] \} \end{aligned}$$

Algorithm NNCS analysis (discrete).

**Input:** NNCS  $(X, U, F, \mathcal{N}, \delta, \mathcal{X}_0)$ , HA  $\mathcal{H}'$ , initial set  $\mathcal{I}_0$ ,  $T = K\delta$ **Output:** Reachable states  $\mathcal{R}$  over [0, T]1: procedure NNCSREACH(I\_0, K)  $\mathcal{R} \leftarrow \text{CONSTRUCTFLOWPIPE}(\mathcal{I}_0, \delta);$ 2.  $\mathcal{X}_i \leftarrow \text{GETRELEVANTSETS}(\mathcal{R}, \delta);$ 3. 4:  $S_{i-1} \leftarrow \{(\mathcal{X}_i, F_0, \mathcal{I}_0)\}$ for i = 1 to K do 5:  $S_i \leftarrow \emptyset$ : 6: for  $(\mathcal{X}_i, F_i, \mathcal{I}_i) \in \mathcal{S}_{i-1}$  do 7:  $(\mathcal{X}', \mathcal{F}', \mathcal{I}') \leftarrow \text{COMPUTESETFLOWMAPS}(\mathcal{X}_i, \mathcal{F}_i, \mathcal{I}_i);$ 8.  $S_i \leftarrow S_i \cup \{(\mathcal{X}', F', \mathcal{I}')\}$ : Q٠ 10: end for for j = 1 to  $|S_j|$  do 11:  $\mathcal{R}_i \leftarrow \text{CONSTRUCTFLOWPIPE}(\mathcal{I}_i, \delta);$ 12:  $\mathcal{R} \leftarrow \mathcal{R} \cup \mathcal{R}_i$ 13: end for 14  $S_{i-1} \leftarrow S_i$ : 15 16  $\mathcal{X}_i \leftarrow \text{GETRELEVANTSETS}(\mathcal{R}, (i+1)\delta);$ end for 17: return  $\mathcal{R}$ 18:

19: end procedure

 $t_4 = 4\delta$ 



$$\mathcal{N}(\Theta(\mathcal{X}_j)).range = \{0, K\rho\}$$
  
$$\implies \mathcal{S}_i = \{(\mathcal{X}_i, f_{1,x}, \mathcal{I}'), (\mathcal{X}_i, f_{2,x}, \mathcal{I}'')\}$$

Algorithm NNCS analysis (discrete).

**Input:** NNCS  $(X, U, F, \mathcal{N}, \delta, \mathcal{X}_0)$ , HA  $\mathcal{H}'$ , initial set  $\mathcal{I}_0$ ,  $T = K\delta$ **Output:** Reachable states  $\mathcal{R}$  over [0, T]1: procedure NNCSREACH(I\_0, K)  $\mathcal{R} \leftarrow \text{CONSTRUCTFLOWPIPE}(\mathcal{I}_0, \delta);$ 2.  $\mathcal{X}_i \leftarrow \text{GETRELEVANTSETS}(\mathcal{R}, \delta);$ 3. 4:  $S_{i-1} \leftarrow \{(\mathcal{X}_i, F_0, \mathcal{I}_0)\}$ for i = 1 to K do 5:  $S_i \leftarrow \emptyset$ : 6: for  $(\mathcal{X}_i, F_i, \mathcal{I}_i) \in \mathcal{S}_{i-1}$  do 7:  $(\mathcal{X}', \mathcal{F}', \mathcal{I}') \leftarrow \text{COMPUTESETFLOWMAPS}(\mathcal{X}_i, \mathcal{F}_i, \mathcal{I}_i);$ 8.  $S_i \leftarrow S_i \cup \{(\mathcal{X}', F', \mathcal{I}')\}$ : Q٠ 10: end for for i = 1 to  $|S_i|$  do 11:  $\mathcal{R}_i \leftarrow \text{CONSTRUCTFLOWPIPE}(\mathcal{I}_i, \delta);$ 12:  $\mathcal{R} \leftarrow \mathcal{R} \cup \mathcal{R}_i$ 13: end for 14  $S_{i-1} \leftarrow S_i$ : 15 16  $\mathcal{X}_i \leftarrow \text{GETRELEVANTSETS}(\mathcal{R}, (i+1)\delta);$ end for 17: return  $\mathcal{R}$ 18:

19: end procedure

 $t_4 = 4\delta$ 



Branching into two flowpipes with different flows
# NNCS Analysis: Pseudo-code

Algorithm NNCS analysis (discrete).

**Input:** NNCS  $(X, U, F, \mathcal{N}, \delta, \mathcal{X}_0)$ , HA  $\mathcal{H}'$ , initial set  $\mathcal{I}_0$ ,  $T = K\delta$ **Output:** Reachable states  $\mathcal{R}$  over [0, T]1: procedure NNCSREACH(I\_0, K)  $\mathcal{R} \leftarrow \text{CONSTRUCTFLOWPIPE}(\mathcal{I}_0, \delta);$ 2.  $\mathcal{X}_i \leftarrow \text{GETRELEVANTSETS}(\mathcal{R}, \delta);$ 3. 4:  $S_{i-1} \leftarrow \{(\mathcal{X}_i, F_0, \mathcal{I}_0)\}$ for i = 1 to K do 5:  $S_i \leftarrow \emptyset$ : 6: for  $(\mathcal{X}_i, F_i, \mathcal{I}_i) \in \mathcal{S}_{i-1}$  do 7:  $(\mathcal{X}', \mathcal{F}', \mathcal{I}') \leftarrow \text{COMPUTESETFLOWMAPS}(\mathcal{X}_i, \mathcal{F}_i, \mathcal{I}_i);$ 8.  $S_i \leftarrow S_i \cup \{(\mathcal{X}', F', \mathcal{I}')\}$ : Q٠ 10: end for for j = 1 to  $|S_j|$  do 11:  $\mathcal{R}_i \leftarrow \text{CONSTRUCTFLOWPIPE}(\mathcal{I}_i, \delta);$ 12:  $\mathcal{R} \leftarrow \mathcal{R} \cup \mathcal{R}_i$ 13: end for 14  $S_{i-1} \leftarrow S_i$ 15 16  $\mathcal{X}_i \leftarrow \text{GETRELEVANTSETS}(\mathcal{R}, (i+1)\delta);$ 17: end for return  $\mathcal{R}$ 18:

19: end procedure

 $t_5 = 5\delta$ 



# NNCS Analysis: Pseudo-code

Algorithm NNCS analysis (discrete).

**Input:** NNCS  $(X, U, F, \mathcal{N}, \delta, \mathcal{X}_0)$ , HA  $\mathcal{H}'$ , initial set  $\mathcal{I}_0$ ,  $T = K\delta$ **Output:** Reachable states  $\mathcal{R}$  over [0, T]1: procedure NNCSREACH(I\_0, K)  $\mathcal{R} \leftarrow \text{CONSTRUCTFLOWPIPE}(\mathcal{I}_0, \delta);$ 2.  $\mathcal{X}_i \leftarrow \text{GETRELEVANTSETS}(\mathcal{R}, \delta);$ 3. 4:  $S_{i-1} \leftarrow \{(\mathcal{X}_i, F_0, \mathcal{I}_0)\}$ for i = 1 to K do 5:  $S_i \leftarrow \emptyset$ : 6: for  $(\mathcal{X}_i, F_i, \mathcal{I}_i) \in \mathcal{S}_{i-1}$  do 7:  $(\mathcal{X}', \mathcal{F}', \mathcal{I}') \leftarrow \text{COMPUTESETFLOWMAPS}(\mathcal{X}_i, \mathcal{F}_i, \mathcal{I}_i);$ 8.  $S_i \leftarrow S_i \cup \{(\mathcal{X}', F', \mathcal{I}')\}$ : Q٠ 10: end for for j = 1 to  $|S_j|$  do 11:  $\mathcal{R}_i \leftarrow \text{CONSTRUCTFLOWPIPE}(\mathcal{I}_i, \delta);$ 12:  $\mathcal{R} \leftarrow \mathcal{R} \cup \mathcal{R}_i$ 13: end for 14  $S_{i-1} \leftarrow S_i$ 15 16  $\mathcal{X}_i \leftarrow \text{GETRELEVANTSETS}(\mathcal{R}, (i+1)\delta);$ end for 17: return  $\mathcal{R}$ 18: 19: end procedure

 $t_6 = 6\delta$ 



Unsafe region entered

# NNCS Analysis: Pseudo-code

Algorithm NNCS analysis (discrete).

**Input:** NNCS  $(X, U, F, \mathcal{N}, \delta, \mathcal{X}_0)$ , HA  $\mathcal{H}'$ , initial set  $\mathcal{I}_0$ ,  $T = K\delta$ **Output:** Reachable states  $\mathcal{R}$  over [0, T]1: procedure NNCSREACH(I\_0, K)  $\mathcal{R} \leftarrow \text{CONSTRUCTFLOWPIPE}(\mathcal{I}_0, \delta);$ 2.  $\mathcal{X}_i \leftarrow \text{GETRELEVANTSETS}(\mathcal{R}, \delta);$ 3. 4:  $S_{i-1} \leftarrow \{(\mathcal{X}_i, F_0, \mathcal{I}_0)\}$ for i = 1 to K do 5:  $S_i \leftarrow \emptyset$ : 6: for  $(\mathcal{X}_i, F_i, \mathcal{I}_i) \in \mathcal{S}_{i-1}$  do 7:  $(\mathcal{X}', F', \mathcal{I}') \leftarrow \text{COMPUTESETFLOWMAPS}(\mathcal{X}_i, F_i, \mathcal{I}_i);$ 8:  $S_i \leftarrow S_i \cup \{(\mathcal{X}', F', \mathcal{I}')\}$ : Q٠ 10: end for for j = 1 to  $|S_j|$  do 11:  $\mathcal{R}_i \leftarrow \text{CONSTRUCTFLOWPIPE}(\mathcal{I}_i, \delta);$ 12:  $\mathcal{R} \leftarrow \mathcal{R} \cup \mathcal{R}_i$ 13: end for 14.  $S_{i-1} \leftarrow S_i$ ; 15 16.  $\mathcal{X}_i \leftarrow \text{GETRELEVANTSETS}(\mathcal{R}, (i+1)\delta);$ 17: end for return  $\mathcal{R}$ 18:

19: end procedure

 $t_{21} = 21\delta$ 



Full time horizon reached

Theorem (Over-approximation of First Segment) Let  $\mathfrak{N} = (X, U, F, \mathcal{N}, \delta, \mathcal{X}_0)$  be an NNCS. Then  $\mathcal{R}_{[0,\delta]}(\mathcal{X}_0) \subseteq \Omega_{[0,\delta]}(\mathcal{X}_0, \mathcal{U}).$  Theorem (Over-approximation of First Segment) Let  $\mathfrak{N} = (X, U, F, \mathcal{N}, \delta, \mathcal{X}_0)$  be an NNCS. Then  $\mathcal{R}_{[0,\delta]}(\mathcal{X}_0) \subseteq \Omega_{[0,\delta]}(\mathcal{X}_0, \mathcal{U}).$ 

$$\implies \mathcal{R}_{[0,T]}(\mathcal{X}_0) \subseteq \Omega_{[0,T]}(\mathcal{X}_0, \mathcal{U}) = \bigcup_{i=1}^{K} \Omega_{[(i-1)\delta, i\delta]}(\mathcal{I}_{i-1}, \mathcal{U}).$$

Theorem (Over-approximation of First Segment) Let  $\mathfrak{N} = (X, U, F, \mathcal{N}, \delta, \mathcal{X}_0)$  be an NNCS. Then  $\mathcal{R}_{[0,\delta]}(\mathcal{X}_0) \subseteq \Omega_{[0,\delta]}(\mathcal{X}_0, \mathcal{U}).$ 

$$\implies \mathcal{R}_{[0,T]}(\mathcal{X}_0) \subseteq \Omega_{[0,T]}(\mathcal{X}_0,\mathcal{U}) = \bigcup_{i=1}^K \Omega_{[(i-1)\delta,i\delta]}(\mathcal{I}_{i-1},\mathcal{U}).$$

#### Theorem

Soundness Let  $\mathfrak{N} = (X, U, F, \mathcal{N}, \delta, \mathcal{X}_0)$  be an NNCS,  $I \subseteq X$  the initial state set,  $\overline{S} \subseteq X$  the set of bad states. Then  $\mathfrak{N}$  is safe iff  $\Omega_{[0, T]}(I) \cap \overline{S} = \emptyset$ .

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Thermostat

- Global time horizon T = 2.1
- Step size  $\delta = 0.1$
- Segment size 0.05
- Initial set *x* ∈ [19.5, 20.5]
- Rod reactor
  - Global time horizon T = 17
  - Step size  $\delta = 1$
  - Segment size 0.2
  - Initial set  $x \in [510, 520], c_1 = c_2 = 20$

## Result on Thermostat



## Rod Reactor Hybrid Automaton



## Result on Rod Reactor



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Run-time

Run 100 times

### Run-time

- Run 100 times
- Accuracy
  - Area of projected flowpipes

### Run-time

- Run 100 times
- Accuracy
  - Area of projected flowpipes
- Complexity
  - Branching and number of flowpipes

Rep	Mean	Median	Min	Max	Std		
Thermostat							
Box	880.73	864.33	657.49	3433.20	279.67		
$\mathcal H$ -polytope	4426.76	4418.52	3759.03	4980.86	186.12		
$\mathcal V$ -polytope	2544.66	2496.54	1650.08	5090.44	355.09		
Rod reactor							
Box	6969.49	6953.31	4877.65	9638.46	474.57		
$\mathcal V$ -polytope	93445.35	91314.90	87657.10	116347.00	4646.71		

Projection	Box	$\mathcal H$ -polytope	$\mathcal V$ -polytope				
Thermostat							
t,x	11.0967	10.4597	6.8272				
Rod Reactor							
t,x	8226.88	-	8087.65				
$t, c_1$	49.02	-	5.82292				

# Accuracy Thermostat



# Accuracy Rod Reactor



# Complexity Thermostat



# Complexity Rod Reactor



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- Flowpipe construction
- Exact star-based neural network reachability analysis

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- Exact star-based neural network reachability analysis
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- Quantified run-time, accuracy, and complexity
- Implemented with HyPro [Schupp et al. 2017]

Refine controller output to reduce over-approximation errors

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- Analyze theoretical time and space complexity
- Overcome numerical challenges

Thank you for your attention!



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# References I

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## Wrapping Effects of Box Representation





(a) Step-size dependent

(b) Inherent over-approximation errors

Wrapping effects in the box representation.

## Step Size Rod Reactor



## Reachability Analysis using Staircase Function

Algorithm Star-based exact reachability analysis for one layer with step function Input Input star set  $I = [\Theta_1 \cdots \Theta_N]$ , ordered classes  $[h_1 \cdots h_m]$  with intervals  $[A_1 \cdots A_m]$ Output Exact reachable set R 1: procedure LAYERREACH(1, W, b)  $\mathcal{R} \leftarrow \emptyset$ : 2. 3 for i = 1 : N do  $I_1 \leftarrow W * \Theta_i + b = \langle Wc_i + b, WV_i, P_i \rangle;$ 4. 5  $\mathcal{R}_1 \leftarrow I_1$ ; for i = 1 : n do 6:  $[l_i, u_i] \leftarrow l_1.range(i);$  $\triangleright$   $l_i \leq x_i \leq u_i, x_i \in l_1[i]$ 7:  $\mathcal{R}_1 \leftarrow \text{STEPSTAIRCASE}(\mathcal{R}_1, i, l_i, u_i);$ 8. Q٠ end for 10:  $\mathcal{R} \leftarrow \mathcal{R} \cup \mathcal{R}_1$ : 11end for 12. return R: 13: end procedure 14: procedure STEPSTAIRCASE(1, i, li, ui)  $\tilde{\mathcal{R}} \leftarrow \emptyset$ : 15  $\tilde{I} = [\tilde{\Theta}_1 \cdots \tilde{\Theta}_k];$ 16  $M \leftarrow [e_1 \ e_2 \ \cdots \ e_{i-1} \ 0 \ e_{i+1} \ \cdots \ e_n];$ 17: Intermediate representations for i = 1 : k do 18:  $\mathcal{R}_1 \leftarrow \emptyset$ ; 19  $\tilde{\Theta}_i = \langle \tilde{c}_i, \tilde{V}_i, \tilde{P}_i, \rangle;$ 20.  $i_{min} \leftarrow \arg\min_{i'} \{h_{i'} \mid I_i \in A_{i'}, i' < m\};$ Index of smallest class 21  $i_{max} \leftarrow \arg \max_{i'} \{ h_{i'} \mid u_i \in A_{i'}, i' \leq m \};$ 22. Index of greatest class 23 for  $i' = i_{min} : i_{max}$  do 24  $v_1 \leftarrow h_{i'} \cdot e_i;$  $\tilde{\Theta}'_{i'} \leftarrow \langle M\tilde{c}_j + v_1, M\tilde{V}_i, \tilde{P}_i, \rangle;$ 25  $\mathcal{R}_1 \leftarrow \mathcal{R}_1 \cup \tilde{\Theta}'_{i'};$ 26: 27. end for  $\tilde{\mathcal{R}} \leftarrow \tilde{\mathcal{R}} \cup \mathcal{R}_1$ : 28. 29 end for return  $\tilde{\mathcal{R}}$ : 30 31: end procedure

```
Gabriela Jiang
```

## Proof.

- $\blacksquare$  Flowpipe construction with  $\bar{\mathcal{R}}_{[0,\delta]}$  is over-approximation
- 1. controller invocation at time δ: reachable set X<sub>i</sub> ⊆ R
  <sub>[0,δ]</sub> with X<sub>i</sub> ∩ R<sub>[δ,δ]</sub> is used to obtain the current state variable valuations for creating the input star Θ<sub>I</sub>, which by design encompasses the reachable set at time δ.
- Exactness of star-based reachability algorithm
- Initial set for the first time step is  $\mathcal{I}_k = e^{\delta A_0} \mathcal{X}_0$ , for  $A_0$  over  $[0,\delta]$ .
- Hence  $\mathcal{I}_k \subseteq \overline{\mathcal{R}}_{[0,\delta]}$ .
- Prediction yields a valid flow such that in the next iteration, flowpipe construction can be applied to the transformed initial set.