Modeling and Analysis of Hybrid Systems

Model checking timed automata

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Informatik 2 - Theory of Hybrid Systems
RWTH Aachen University

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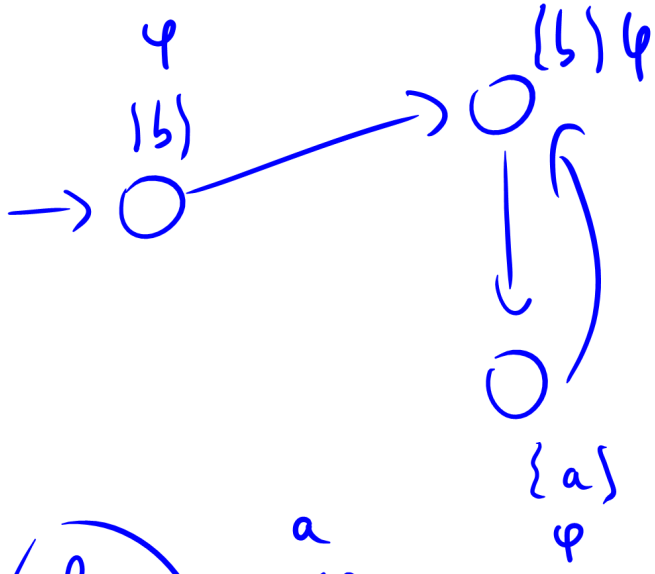
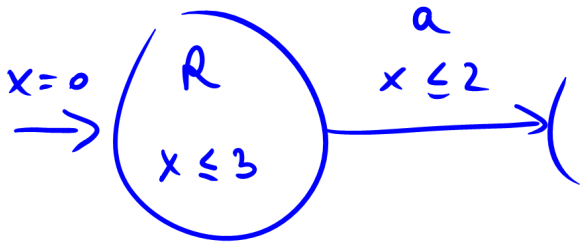
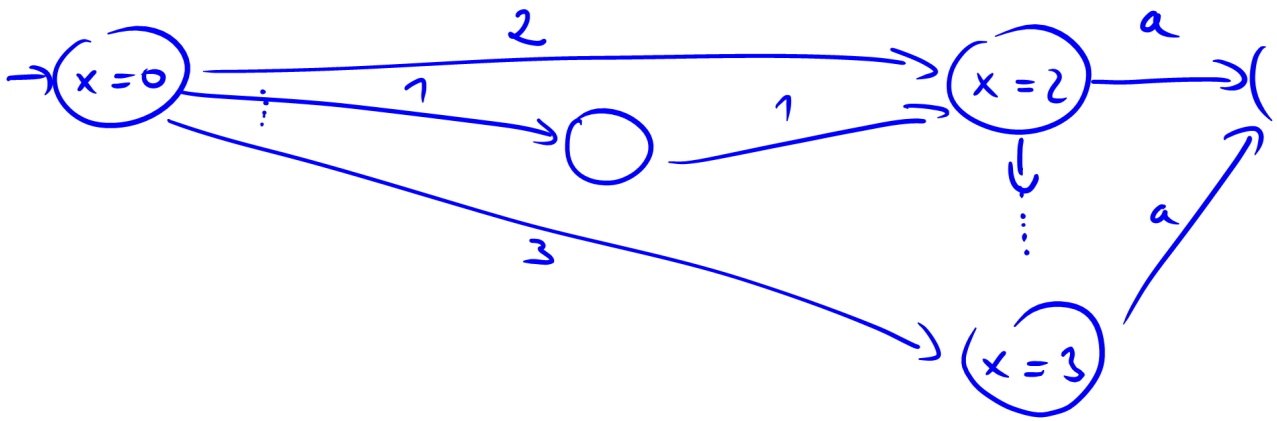


CTL:

$\varphi = A(b \; U \; a)$

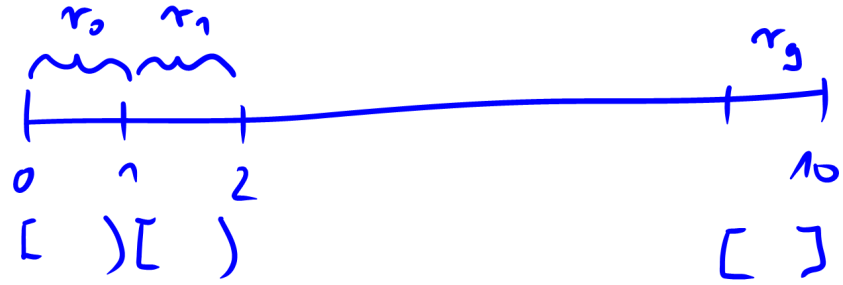




T: $x \in [0,10] \subseteq \mathbb{R}_{\geq 0}$ ✓

"VI" ⇓

K: $x \in [0,10] \subseteq \mathbb{Z}$

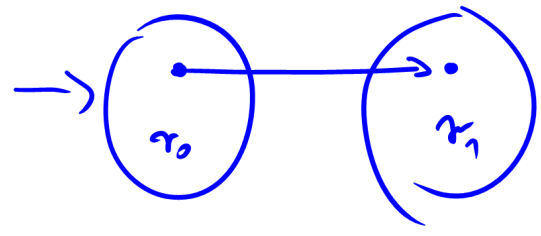


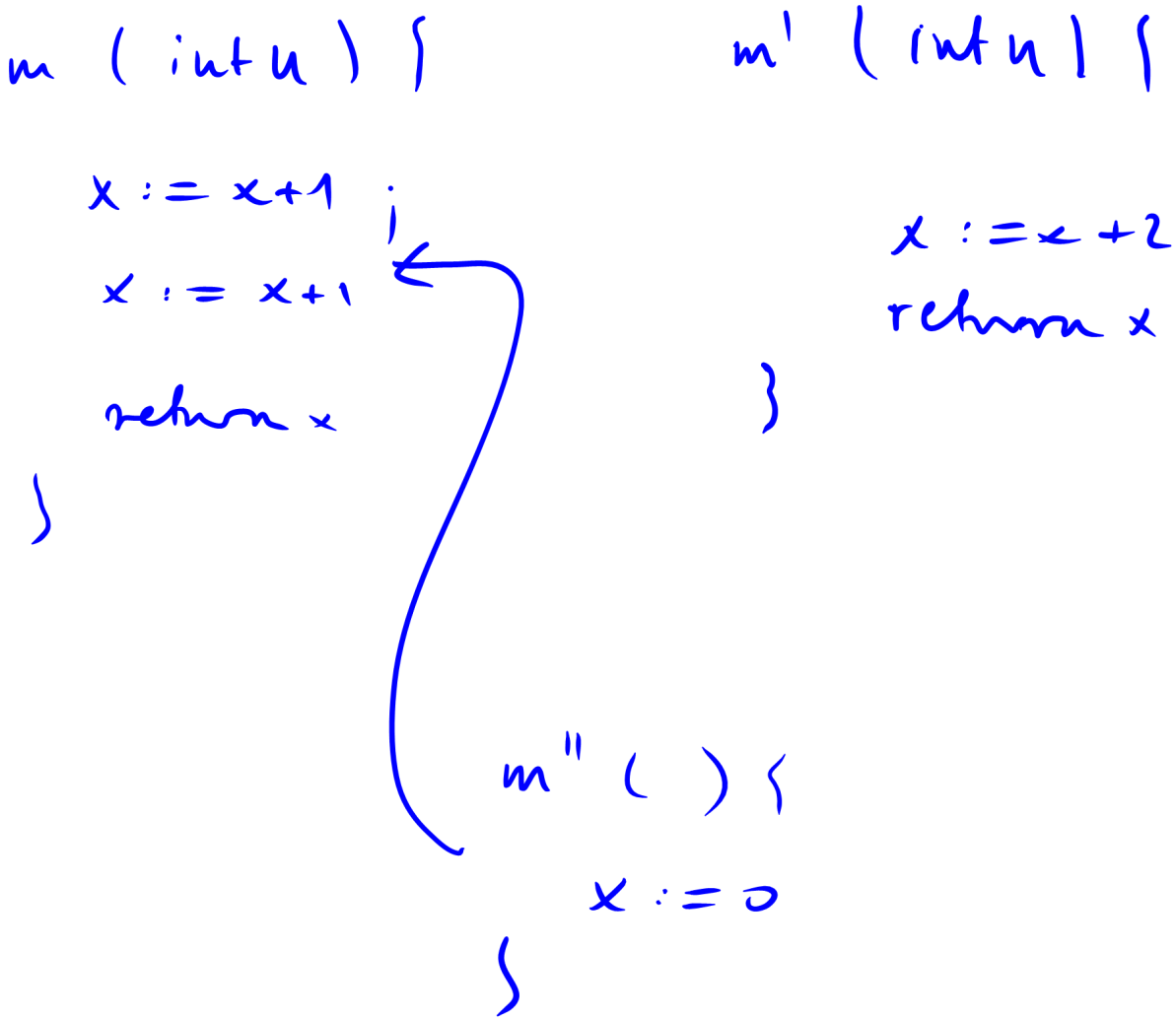
$K \models \exists \psi \Rightarrow T \models \exists \psi$

$K \not\models \forall \psi \Rightarrow T \not\models \forall \psi$ → r_0 → r_1

$K \models \forall \psi \Rightarrow x$

$K \not\models \exists \psi \Rightarrow x$




m (int x) {
x := x+1 ;
x := x+1
return x
}
m' (int x) {
x := x+2
return x
}
m'' ( ) {
x := 0
}

Basic method: Abstraction

- Given: a concrete system
  (here: a timed automaton)
- Goal: reduce the size of the system by abstraction
  (here: reduce the infinite state space to a finite one)
- Result: abstract system
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- Two systems P and P' have the same observable behaviour iff for each context C we have that $[\![C[P]]\!] = [\![C[P']]\!]$.
  ($C[P]$: the composition of C and P, $[\![\cdot]\!]$: (global) semantics)
- E.g., for programs it could mean the same input-output behaviour. For model checking we require that they satisfy the same formulas of the underlying logic.
  (here: TCTL)

TCTL model checking

Input: timed automaton $\mathcal{T}$, TCTL formula ψ
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1. Eliminate the timing parameters from the TCTL formula ψ, resulting in a CTL formula $\hat{\psi}$ (with clock constraints as atomic propositions).
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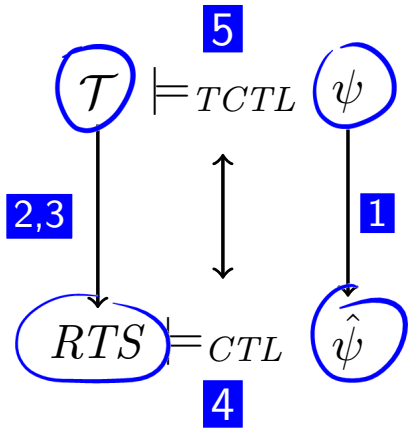

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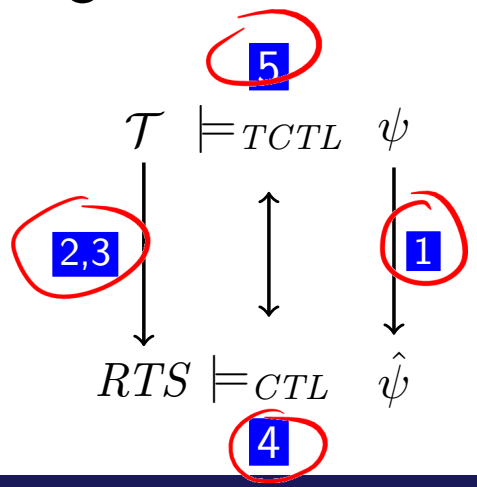

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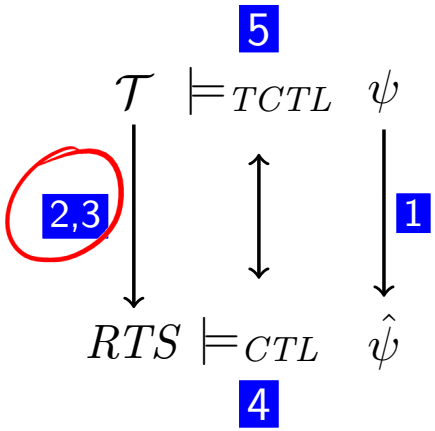

Keywords:

Finite abstraction

Equivalence relation, equivalence classes

Bisimulation

And what does it mean in our context?

2. Finite state space abstraction

$\mathcal{T}, \psi \quad \sigma \cong \sigma' \Rightarrow (\sigma \models \psi \Leftrightarrow \sigma' \models \psi)$

We search for an equivalence relation $\cong$ on states, such that equivalent states satisfy the same (sub)formulae ψ' occurring in the timed automaton $\mathcal{T}$ or in the specification ψ:

$$\sigma \cong \sigma' \quad \Rightarrow \quad (\sigma \models \psi' \quad \textit{iff} \quad \sigma' \models \psi') \,.$$

Since the set of such (sub)formulae is finite, we strive for a finite number of equivalence classes.

Bisimulation on the state space of an LSTS

Definition

Let $LSTS = (\Sigma, Lab, Edge, Init)$ be a state transition system, AP a set of atomic propositions, and $L : \Sigma \to 2^{AP}$ a labeling function over AP.
A bisimulation for $LSTS$ is an equivalence relation $\approx \subseteq \Sigma \times \Sigma$ such that for all $\sigma_1 \approx \sigma_2$

1. $L(\sigma_1) = L(\sigma_2)$
2. for all $\sigma_1' \in \Sigma$ with $\sigma_1 \xrightarrow{a} \sigma_1'$ there exists $\sigma_2' \in \Sigma$ such that $\sigma_2 \xrightarrow{a} \sigma_2'$ and $\sigma_1' \approx \sigma_2'$.

Time abstract bisimulation

Definition

Let $\mathcal{T} = (Loc, \mathcal{C}, Lab, Edge, Inv, Init)$ be a timed automaton, AP a set of atomic propositions, and $L : \Sigma \to 2^{AP}$.

A time abstract bisimulation on $\mathcal{T}$ is an equivalence relation $\approx \subseteq \Sigma \times \Sigma$ such that for all $\sigma_1, \sigma_2 \in \Sigma$ satisfying $\sigma_1 \approx \sigma_2$

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- for all $\sigma_1' \in \Sigma$ with $\sigma_1 \xrightarrow{t_1} \sigma_1'$ there is a $\sigma_2' \in \Sigma$ such that $\sigma_2 \xrightarrow{t_2} \sigma_2'$ and $\sigma_1' \approx \sigma_2'$.

Bisimulation

Lemma

Assume a timed automaton $\mathcal{T}$ with state space Σ, and a time-abstract bisimulation $\approx \subseteq \Sigma \times \Sigma$ on $\mathcal{T}$.
Then for all $\sigma, \sigma' \in \Sigma$ with $\sigma \approx \sigma'$ we have that for each path

$$\pi : \sigma \xrightarrow{\alpha_1} \sigma_1 \xrightarrow{\alpha_2} \sigma_2 \xrightarrow{\alpha_3} \dots$$

of $\mathcal{T}$ there exists a path

$$\pi' : \sigma' \xrightarrow{\alpha'_1} \sigma'_1 \xrightarrow{\alpha'_2} \sigma'_2 \xrightarrow{\alpha'_3} \dots$$

of $\mathcal{T}$ such that for all i

- $\sigma_i \approx \sigma'_i$,
- $\alpha_i = \alpha'_i$ *if* $\alpha_i \in Lab$ *and*
- $\alpha_i, \alpha'_i \in \mathbb{R}_{\geq 0}$ *otherwise.*

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Since, in general,

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different locations in $\mathcal{T}$ are different, only states (l, ν) and (l', ν') satisfying $l = l'$ should be equivalent.

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Notation:

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I.e., only states (l, ν) and (l, ν') satisfying

$$\lfloor \nu(x) \rfloor = \lfloor \nu'(x) \rfloor \; and \; frac(\nu(x)) = 0 \; iff \; frac(\nu'(x)) = 0$$

for all $x \in \mathcal{C}$ should be equivalent.

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Let c_x be the largest constant which a clock x is compared to in $\mathcal{T}$ or in ψ. Then there is no observation which could distinguish between the x-values in (l, ν) and (l, ν') if $\nu(x) > c_x$ and $\nu'(x) > c_x$.

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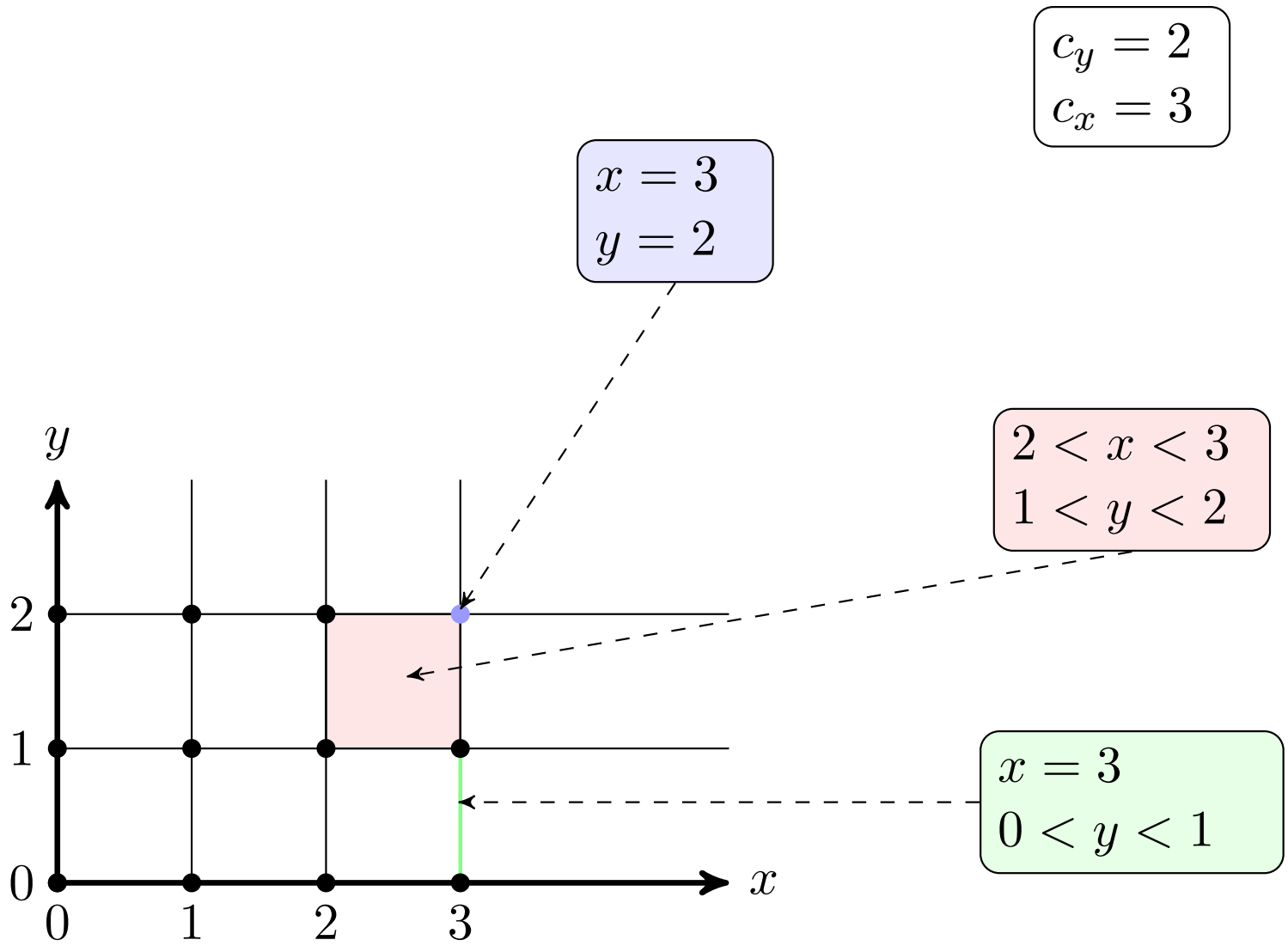

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As the following example illustrates, we must make a further refinement of the abstraction, since it does not distinguish between states satisfying different formulae.

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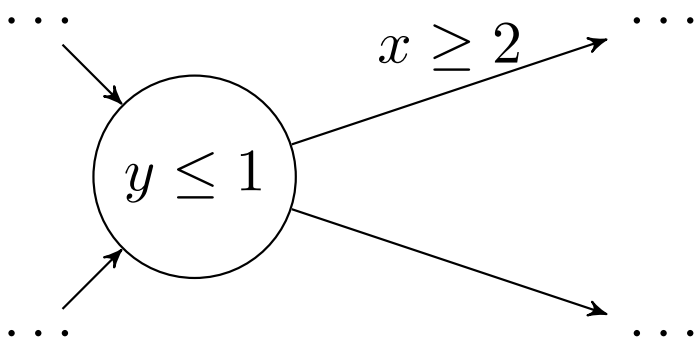

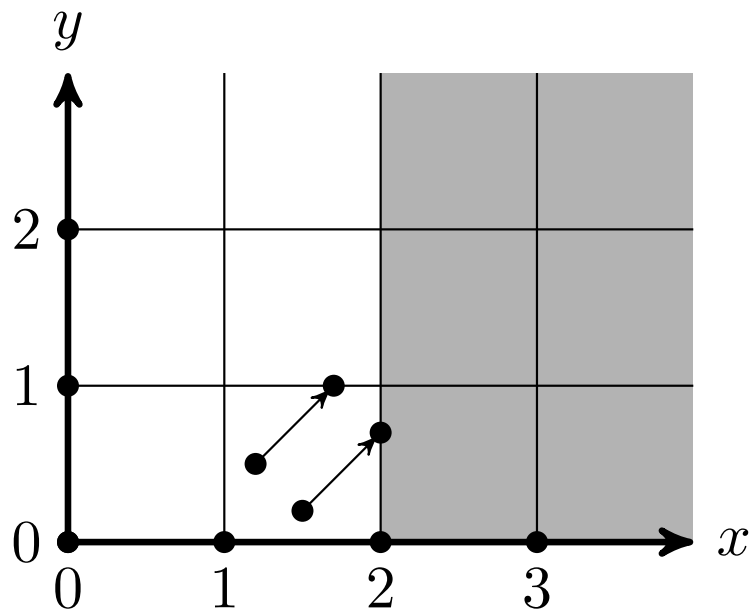

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What we need is a refinement taking the order of the fractional parts of the clock values into account. However, again only for values below the largest constants to which the clocks get compared.
I.e., only states (l, ν) and (l, ν') satisfying

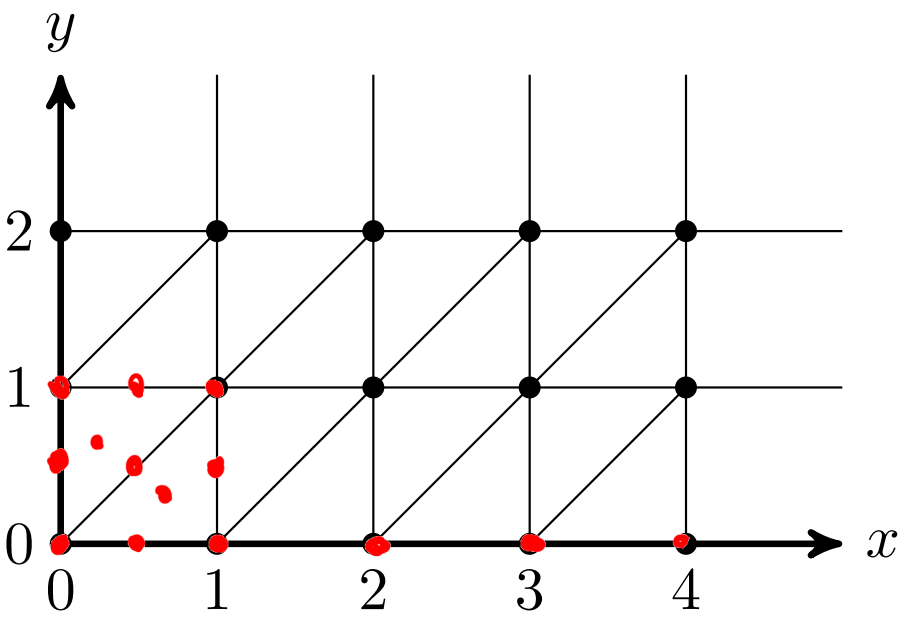
$$
\begin{aligned}
&(\nu(x), \nu'(x) > c_x \wedge \nu(y), \nu'(y) > c_x) \quad \vee \\
&\quad (\quad frac(\nu(x)) < frac(\nu(y)) \quad iff \quad frac(\nu'(x)) < frac(\nu'(y)) \quad \wedge \\
&\qquad frac(\nu(x)) = frac(\nu(y)) \quad iff \quad frac(\nu'(x)) = frac(\nu'(y)) \quad \wedge \\
&\qquad frac(\nu(x)) > frac(\nu(y)) \quad iff \quad frac(\nu'(x)) > frac(\nu'(y)))
\end{aligned}
$$

for all $x, y \in \mathcal{C}$ should be equivalent.
Because of symmetry the following is also sufficient:

$$
\begin{aligned}
&(\nu(x), \nu'(x) > c_x \wedge \nu(y), \nu'(y) > c_y) \quad \vee \\
&\quad (frac(\nu(x)) \leq frac(\nu(y)) \quad iff \quad frac(\nu'(x)) \leq frac(\nu'(y)))
\end{aligned}
$$

for all $x, y \in \mathcal{C}$ should be equivalent.

2. Finite state space abstraction



$$c_y = 2$$
$$c_x = 4$$

finite index

Definition

For a timed automaton $\mathcal{T}$ and a TCTL formula ψ, both over a clock set $\mathcal{C}$, we define the clock equivalence relation $\cong \subseteq \Sigma \times \Sigma$ by $(l, \nu) \cong (l', \nu')$ iff $l = l'$ and

- for all $x \in \mathcal{C}$, either $\nu(x) > c_x \wedge \nu'(x) > c_x$ or

$$\lfloor \nu(x) \rfloor = \lfloor \nu'(x) \rfloor \ \wedge \ (frac(\nu(x)) = 0 \quad iff \quad frac(\nu'(x)) = 0)$$

- for all $x, y \in \mathcal{C}$ if $\nu(x), \nu'(x) \leq c_x$ and $\nu(y), \nu'(y) \leq c_y$ then

$$frac(\nu(x)) \leq frac(\nu(y)) \quad iff \quad frac(\nu'(x)) \leq frac(\nu'(y)).$$

The clock region of an evaluation $\nu \in V$ is the set $[\nu] = \{\nu' \in V \mid \nu \cong \nu'\}$.
The clock region of a state $\sigma = (l, \nu) \in \Sigma$ is the set
$[\sigma] = \{(l, \nu') \in \Sigma \mid \nu \cong \nu'\}$.

2. Finite state space abstraction

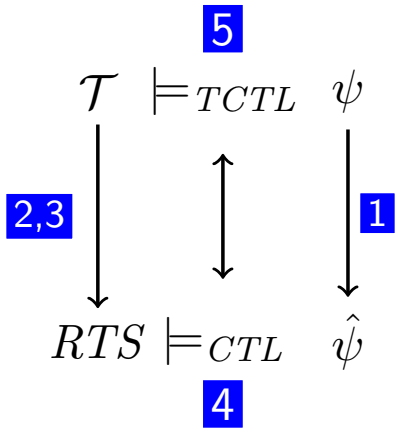
Lemma

Clock equivalence is a bisimulation over $AP' = AP \cup ACC(\mathcal{T}) \cup ACC(\psi)$.

TCTL model checking

Input: timed automaton $\mathcal{T}$, TCTL formula ψ
Output: the answer to the question whether $\mathcal{T} \models \psi$

1. Eliminate the timing parameters from the TCTL formula ψ, resulting in a CTL formula $\hat{\psi}$ (with clock constraints as atomic propositions).
2. Make a finite abstraction of the state space of $\mathcal{T}$.
3. Construct abstract state transition system RTS (with the states labeled by clock constraints as atomic propositions) such that $\mathcal{T} \models_{TCTL} \psi$ iff $RTS \models_{CTL} \hat{\psi}$.
4. Apply CTL model checking to check whether $RTS \models_{CTL} \hat{\psi}$.
5. Return the model checking result.



3. The abstract transition system

We have defined regions as abstract states,
now we connect them by abstract transitions.

Two kinds of transitions:
time and discrete

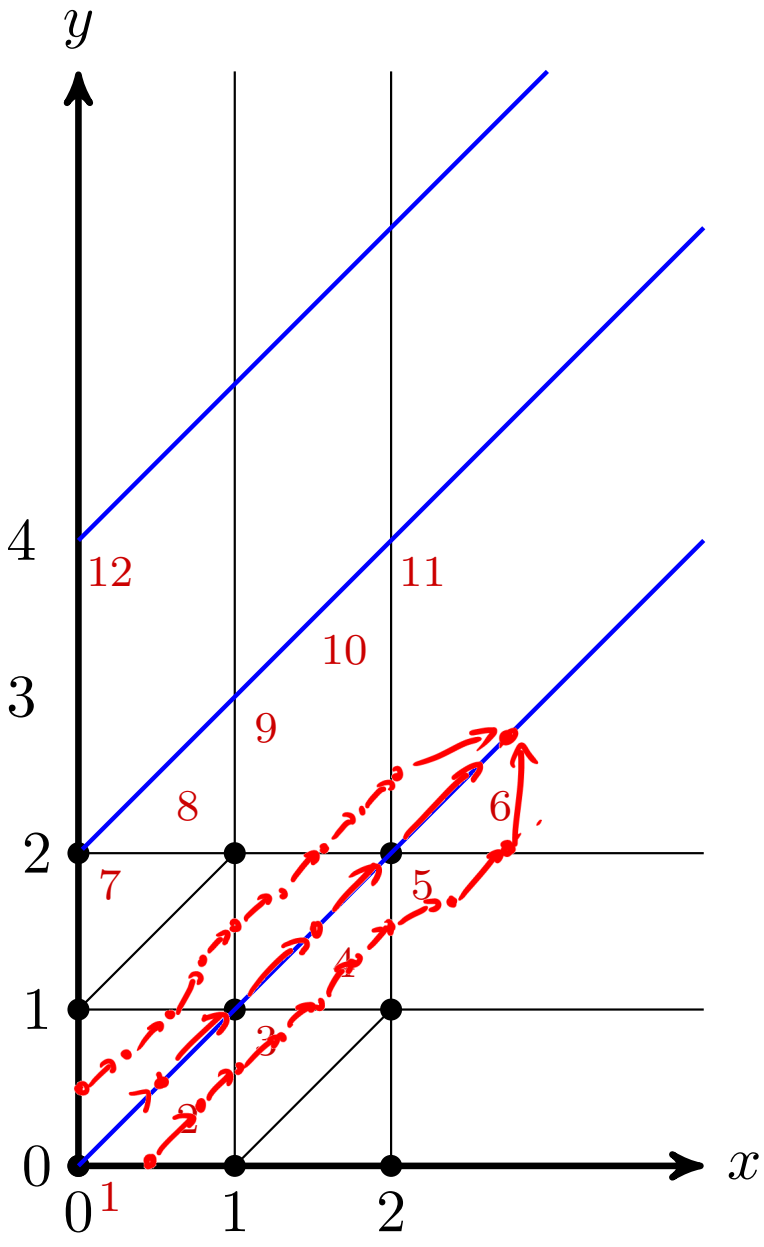
Definition

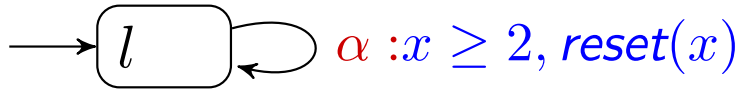
The clock region $r_\infty = \{\nu \in V \mid \forall x \in \mathcal{C}.\ \nu(x) > c_x\}$ is called unbounded. Let r, r' be two clock regions. The region r' is the successor clock region of r, denoted by $r' = succ(r)$, if either

- $r = r' = r_\infty$, or
- $r \neq r_\infty$, $r \neq r'$, and for all $\nu \in r$:

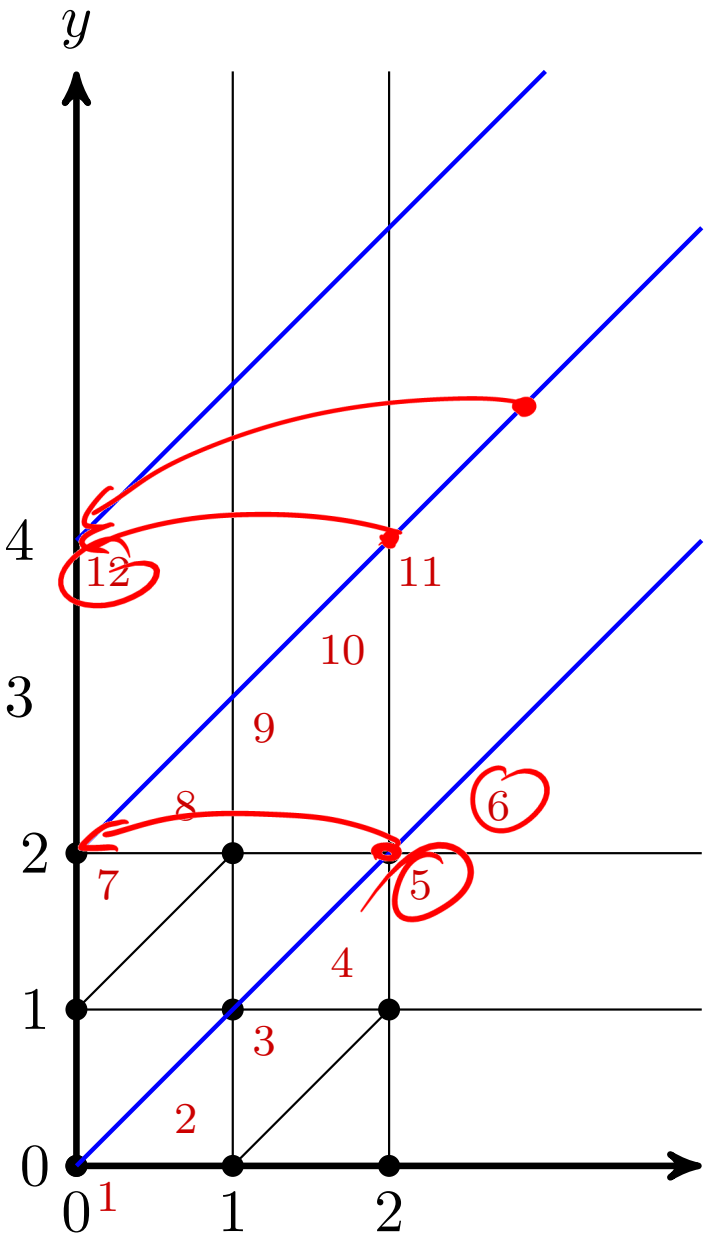
$$\exists d \in \mathbb{R}_{>0}.\ (\nu + d \in r' \ \wedge\ \forall 0 \leq d' \leq d.\ \nu + d' \in (r \cup r')).$$

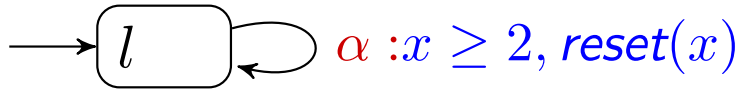
The successor state region is defined as $succ((l, r)) = (l, succ(r))$.



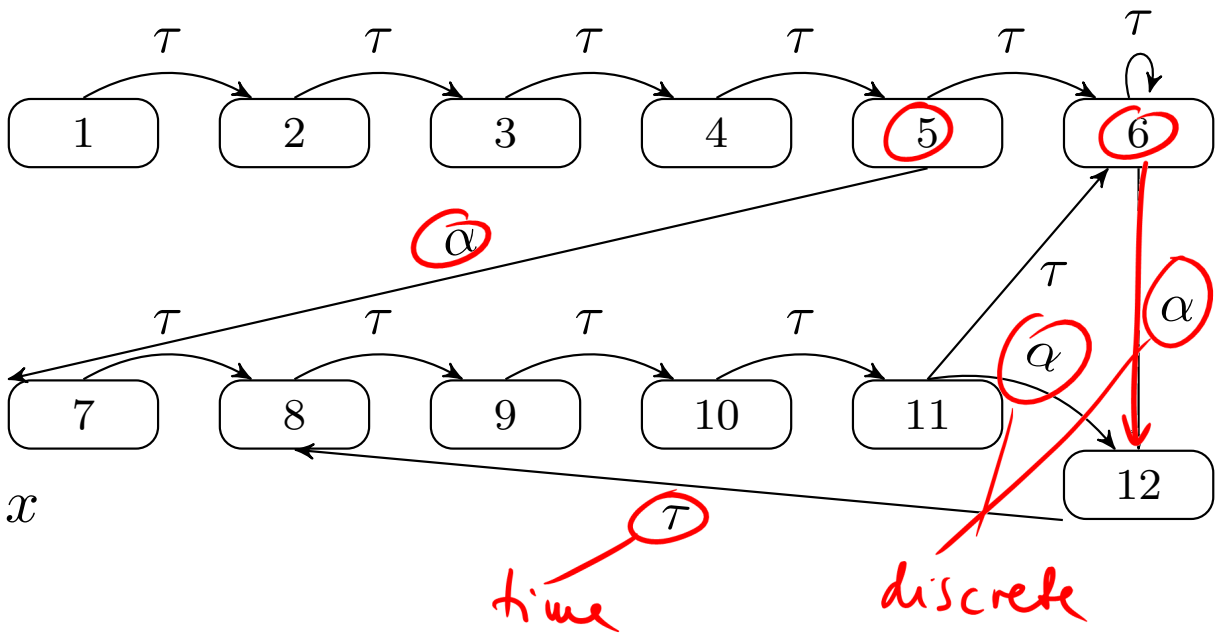


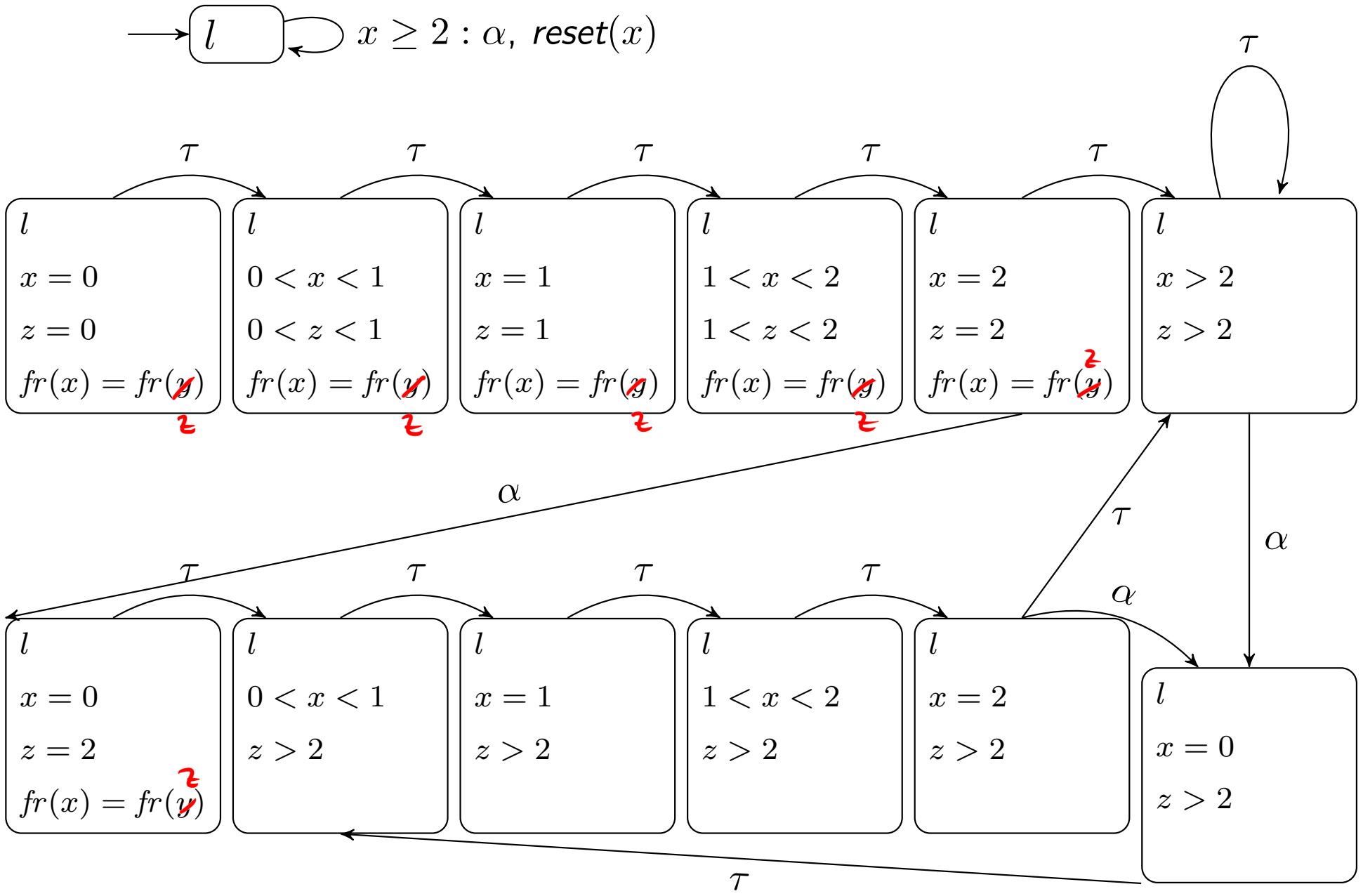
$\mathbf{E}\mathcal{F}^{(0,2]}\ x = 0$





$$\mathbf{E}\mathcal{F}^{(0,2]}\ x = 0$$




l
$x \geq 2 : \alpha,\ reset(x)$
τ
l
$x = 0$
$z = 0$
$fr(x) = fr(y)$
l
$0 < x < 1$
$0 < z < 1$
$fr(x) = fr(y)$
l
$x = 1$
$z = 1$
$fr(x) = fr(y)$
l
$1 < x < 2$
$1 < z < 2$
$fr(x) = fr(y)$
l
$x = 2$
$z = 2$
$fr(x) = fr(y)$
l
$x > 2$
$z > 2$
α
l
$x = 0$
$z = 2$
$fr(x) = fr(y)$
l
$0 < x < 1$
$z > 2$
l
$x = 1$
$z > 2$
l
$1 < x < 2$
$z > 2$
l
$x = 2$
$z > 2$
l
$x = 0$
$z > 2$

3. The abstract transition system

Definition

Let $\mathcal{T} = (Loc, \mathcal{C}, Lab, Edge, Inv, Init)$ be a non-zeno timelock-free timed automaton with an atomic proposition set AP and a labeling function L, and let $\hat{\psi}$ be an unbounded TCTL formula over $\mathcal{C}$ and AP.
The region transition system of $\mathcal{T}$ for $\hat{\psi}$ is a labelled state transition system $\mathcal{RTS}(\mathcal{T}, \psi) = (\Sigma', Lab', Edge', Init')$ with atomic propositions AP' and a labeling function L' such that

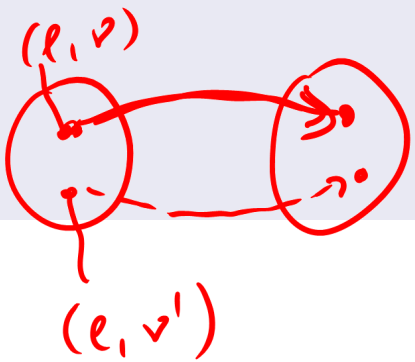
- $\Sigma' = \{(l, [\nu]) \mid (l, \nu) \in \Sigma \wedge \nu \in Inv(l)\}$
- $Init' = \{(l, [\nu]) \in \Sigma' \mid (l, \nu) \in Init\}$
- $AP' = AP \cup ACC(\mathcal{T}) \cup ACC(\hat{\psi})$
- $L'((l, [\nu])) = L(l) \cup \{g \in AP' \backslash AP \mid \nu \models g\}$

and

Definition

$$\frac{(l,\nu) \xrightarrow{a} (l',\nu')}{(l,[\nu]) \xrightarrow{a} (l',[\nu'])} \quad \texttt{Rule}_{\texttt{Discrete}}$$

$$\frac{succ(r) \models Inv(l)}{(l,r) \xrightarrow{t} (l, succ(r))} \quad \texttt{Rule}_{\texttt{Time}}$$



3. The abstract transition system

Lemma

For non-zeno $\mathcal{T}$ and $\pi = s_0 \to s_1 \to \ldots$ an initial, infinite path of $\mathcal{T}$:

- *if π is time-convergent, then there is an index j and a state region (l, r) such that $s_i \in (l, r)$ for all $i \geq j$.*
- *if there is a state region (l, r) with $r \neq r_\infty$ and an index j such that $s_i \in (l, r)$ for all $i \geq j$ then π is time-convergent.*

Lemma

For a non-zeno timed automaton $\mathcal{T}$ and a TCLT formula ψ:

$$\mathcal{T} \models_{TCTL} \psi \quad \textit{iff} \quad RTS(\mathcal{T}, \hat{\psi}) \models_{CTL} \hat{\psi}$$

TCTL model checking

Input: timed automaton $\mathcal{T}$, TCTL formula ψ

Output: the answer to the question whether $\mathcal{T} \models \psi$

1. Eliminate the timing parameters from the TCTL formula ψ, resulting in a CTL formula $\hat{\psi}$ (with clock constraints as atomic propositions).
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4. Apply CTL model checking to check whether $RTS \models_{CTL} \hat{\psi}$.
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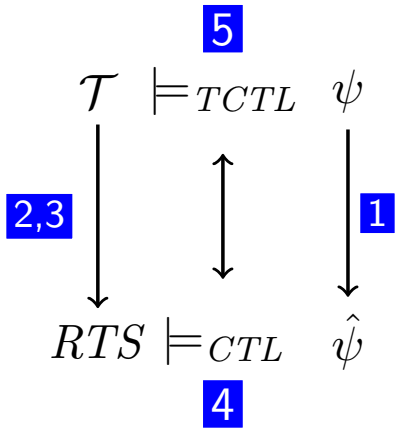

TCTL model checking

The procedure is quite similar to CTL model checking for finite automata.

One difference:

- Handling nested time bounds in TCTL formulae

$$A\ a\ \mathcal{U}^{\leq 2}\ (E\ G\ p)$$

TCTL model checking

Input: timed automaton $\mathcal{T}$, TCTL formula ψ

Output: the answer to the question whether $\mathcal{T} \models \psi$

1. Eliminate the timing parameters from the TCTL formula ψ, resulting in a CTL formula $\hat{\psi}$ (with clock constraints as atomic propositions).
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